

THE COMPLEXITY OF A CLASS OF INFINITE GRAPHS

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Given a class of graphs \mathcal{G} , we say that \mathcal{G} has a *universal* element $G_0 \in \mathcal{G}$ if any other graph $G \in \mathcal{G}$ is isomorphic to a (not necessarily induced) subgraph of G_0 .

In a recent paper (Komjáth-Pach,1989) we have extended this definition, as follows. Let $c(\mathcal{G})$, the *complexity* of a class of graphs \mathcal{G} , be defined as the least cardinality of a subset $\mathcal{G}_0 \subseteq \mathcal{G}$ with the property that any element of \mathcal{G} is isomorphic to a subgraph of some $G_0 \in \mathcal{G}_0$. Obviously, \mathcal{G} has a universal element if and only if $c(\mathcal{G}) = 1$.

In all interesting examples \mathcal{G} is *closed under containment*, i.e., $G \in \mathcal{G}$ implies that $G' \in \mathcal{G}$ for any $G' \subseteq G$. This condition is satisfied e.g. for all classes of graphs which can be obtained in the following way. Given a cardinal κ and a family \mathcal{H} of so-called *forbidden subgraphs*, let $\mathcal{G}_\kappa(\mathcal{H})$ be defined as the class of all graphs with at most κ vertices containing no subgraph isomorphic to any element of \mathcal{H} .

Let \mathcal{G}_k denote the class of all countable graphs containing no k vertex-disjoint cycles. That is, using the above notation, $\mathcal{G}_k = \mathcal{G}_\omega(\mathcal{H}_k)$, where \mathcal{H}_k stands for the family of all (finite) graphs consisting of k vertex-disjoint cycles. In particular, \mathcal{G}_1 is the class of all countable forests. Since \mathcal{G}_k has continuum many elements, its complexity is at most 2^ω .

Theorem 1. *Let $1 < k < \omega$, and let \mathcal{G}_k be the class of all countable graphs containing no k vertex-disjoint cycles. Then $c(\mathcal{G}_k) = \omega$.*

Proof. First we show that $c(\mathcal{G}_k) \leq \omega$.

Let G be a fixed countable graph without k vertex-disjoint cycles. Fix a finite subset $K \subseteq V(G)$ with the property that any cycle of G contains at least one element of K . (Note that the vertex set of any maximal system of cycles in G obviously satisfies this condition. Moreover, by a result of Erdős and Pósa, K can always be chosen so as to have fewer than $ck \log k$ elements.) We will refer to K , as the *kernel* of G . The vertices of G outside the kernel are called *external*. The external vertices

induce a forest in G , and they can be classified according to which elements of K they are connected to. We color two external points with the same *color*, if and only if their sets of neighbors in K are the same. Thus, we obtain a coloring function $\phi: (V(G) - K) \rightarrow \Gamma$, where Γ is the set of colors and $|\Gamma| = 2^{|K|}$.

Let $\mathcal{P} = \{P_1, P_2, \dots, P_m\}$ be a set of at most $|K|$ vertex-disjoint paths in $G-K$, and let v_i and v'_i denote the endpoints of P_i . (We do not exclude the possibility that $v_i = v'_i$, i.e., P_i consists of a single vertex.) The *type of P_i* (with respect to K) is defined by the colors of its endpoints:

$$\text{type}(P_i) = \begin{cases} [\phi(v_i)] & \text{if } v_i = v'_i, \\ [\phi(v_i), \phi(v'_i)] & \text{if } v_i \neq v'_i; \end{cases}$$

$$\text{Type}(\mathcal{P}) = [\text{type}(P_i): 1 \leq i \leq m].$$

(We write $[\cdot]$ instead of $\{\cdot\}$ to indicate that some of the elements may be repeated, i.e., they form a *multiset*.) Furthermore, let $\mathcal{T}(G)$ be defined as the set of all $\text{Type}(\mathcal{P})$, where \mathcal{P} is a system of at most $|K|$ vertex-disjoint paths in $G-K$. Clearly, $\mathcal{T}(G)$ is closed under containment, i.e., $T \in \mathcal{G}(G)$ implies that every submultiset $T' \subseteq T$ also belongs to $\mathcal{T}(G)$.

For each member $T \in \mathcal{T}(G)$, fix a system \mathcal{P}^T of vertex-disjoint paths with $\text{Type}(\mathcal{P}^T) = T$. Put

$$L = \bigcup_{T \in \mathcal{T}(G)} V(\mathcal{P}^T).$$

Evidently, L is a finite set and $L \cap K = \emptyset$. Since $G-K$ is a forest, there are only finitely many external vertices lying on some path connecting two elements of L . Let \bar{L} denote the set obtained from L by adding all of these vertices.

Now $G - (K \cup \bar{L})$ falls into countably many connected components (trees) G_i ($i = 1, 2, \dots$). Every G_i has at most one point adjacent to some element of \bar{L} . If such a point exists, then it is called the *root of G_i* . By the definition of \bar{L} , the root of G_i has only one neighbor in \bar{L} .

Next we describe a procedure that will enable us to add new vertices to G , without creating k vertex-disjoint cycles.

Let G_1 be a fixed component of $G - (K \cup \bar{L})$, and pick a color $\gamma \in \Gamma$ which occurs among the vertices of G_1 at least twice. Thus, one can find two distinct points

$u_1, u_2 \in V(G_1)$ which are adjacent to the same elements of K . Let G' denote the graph obtained from G by adding a new vertex u of color γ and connecting it to any point $w \in V(G_1)$. That is,

$$\begin{aligned} V(G') &= V(G) \cup \{u\} , \\ E(G') &= E(G) \cup \{uv: v \in K , u_1v \in E(G)\} \cup \{uw\} . \end{aligned}$$

Obviously, any cycle of G' passes through at least one element of K . In other words, K is also a kernel of G' , hence it can be used to define $\mathcal{T}(G')$.

Lemma 1. $\mathcal{T}(G') = \mathcal{T}(G)$.

We have to prove only that $\mathcal{T}(G') \subseteq \mathcal{T}(G)$. Assume, in order to obtain a contradiction, that there is a system $\mathcal{P} = \{P_i: 1 \leq i \leq m\}$ of at most $|K|$ vertex-disjoint paths in $G' - K$ such that

$$T = \text{Type}(\mathcal{P}) = [\text{type}(P_i): 1 \leq i \leq m] \in \mathcal{T}(G') - \mathcal{T}(G) .$$

Suppose, without loss of generality, that \mathcal{P} is a *minimal* system satisfying this condition, i.e.,

$$T_j = [\text{type}(P_i): 1 \leq i \leq m , i \neq j] \in \mathcal{T}(G)$$

for every j ($1 \leq j \leq m$).

Clearly, one of the paths P_i (say, P_1) must contain the new vertex u , otherwise $T = \text{Type}(\mathcal{P}) \in \mathcal{T}(G)$. Moreover, u must be an endpoint of P_1 , because the degree of u in $G' - K$ is 1. Let v_i and v'_i denote the (not necessarily distinct) endpoints of P_i . Thus, we can assume that $u = v_1$.

Let $G'_1 \subseteq G'$ denote the tree obtained from G_1 by adding the vertex u and the edge uw .

Observe that no path P_j can be entirely contained in G'_1 . To see this, recall that there is a system \mathcal{P}^{T_j} of vertex-disjoint paths in L with $\text{Type}(\mathcal{P}^{T_j}) = T_j$. So, if P_j were in G'_1 for some $j \neq 1$, then $\mathcal{P}^{T_j} \cup \{P_j\}$ would form a system of vertex-disjoint paths in $G-K$, whose type is T . If $P_1 \subseteq G'_1$, then consider the uniquely determined paths P_{11} and $P_{12} \subseteq G_1$ connecting v'_1 to u_1 and u_2 , respectively. At least one of them (say, P_{11}) is of the same type as P_1 . Hence, $\mathcal{P}^{T_1} \cup \{P_{11}\}$ is a system of vertex-disjoint

paths in $G-K$, whose type is T . In both cases we can conclude that $T \in \mathcal{T}(G)$, contradiction.

Thus, we can assume that v'_1 is not in G'_1 . This implies that G_1 has a root r , and P_1 must pass through r . Let P_{11} denote the (unique) path connecting v'_1 and u_1 in $G-K$. Clearly, P_{11} also passes through r and $\text{type}(P_{11}) = \text{type}(P_1)$. Notice that P_{11} is disjoint from any P_j ($2 \leq j \leq m$), otherwise P_j would lie entirely in G_1 , contradicting our previous observation. Hence, $\{P_{11}, P_2, P_3, \dots, P_m\}$ is a system of vertex-disjoint paths in $G-K$, whose type is T , which is again a contradiction. This completes the proof of Lemma 1.

Lemma 2. *G' has no k vertex-disjoint cycles.*

Assume, for contradiction, that there is a system $\{C_i: 1 \leq i \leq k\}$ of k vertex-disjoint cycles in G' . Since every cycle must visit K , the pieces of the C_i lying outside K form a system \mathcal{P}' of at most $|K|$ vertex-disjoint paths in $G'-K$. By Lemma 1, there exists a system \mathcal{P} of vertex-disjoint paths in $G-K$ such that $\text{Type}(\mathcal{P}) = \text{Type}(\mathcal{P}')$. For every cycle C_i , replace each piece lying outside K by the corresponding path in \mathcal{P} . Thus, we obtain k vertex-disjoint cycles in G , the desired contradiction establishing Lemma 2.

By the repeated application of the above procedure, we can add countably many new vertices to G , to obtain a graph G^* satisfying the conditions summarized in the following statement.

Lemma 3. *Let G be a countable graph without k vertex-disjoint cycles, and let $K, \bar{L} \subseteq V(G)$ be finite sets, as defined above.*

Then there exists a countable graph G^ with the following properties.*

- (i) G^* contains G as an induced subgraph;
- (ii) G^* has no k vertex-disjoint cycles;
- (iii) every cycle of G^* meets K .

Furthermore, let $\phi^: (V(G^*) - K) \rightarrow \Gamma$ be a coloring assigning the same color to two vertices if and only if they are connected to the same elements of K . Let G_i^* ($i = 1, 2, \dots$) denote the connected components of $G^* - (K \cup \bar{L})$.*

- (iv) *Each component G_i^* is connected to \bar{L} by at most one edge;*

(v) if $\gamma \in \Gamma$ is any color assigned to at least two points of G_i^* , then every vertex of G_i^* has infinitely many neighbors of color γ ($i = 1, 2, \dots$).

Let \mathcal{G}_k^* be the family of all countable graphs that can be obtained as G^* for some $G \in \mathcal{G}_k$. Obviously, $\mathcal{G}_k^* \subseteq \mathcal{G}_k$ and every element of \mathcal{G}_k can be embedded into some element of \mathcal{G}_k^* as an induced subgraph. On the other hand, \mathcal{G}_k^* is clearly a *countable* family of graphs. To see this, we have to note only that

- (a) there are only countably many different graphs that can be obtained as the restriction of some $G^* \in \mathcal{G}_k^*$ to the corresponding subset $K \cup \bar{L}$ (because K and \bar{L} are finite);
- (b) there are only countably many different colored graphs (trees) that can be obtained as G_i^* for some $G^* \in \mathcal{G}_k^*$ (because G_i^* is either finite or it is a tree whose every vertex has degree ω , and those points whose color does not appear anywhere else in G_i^* can be situated in this tree in countably many different ways);
- (c) given K and \bar{L} , there are only countably many different ways that a colored tree G_i^* can be connected to these sets (because of Lemma 3 (iv)).

Hence, $c(\mathcal{G}_k) \leq \omega$.

Next we show that $c(\mathcal{G}_k) \geq \omega$. Let K_4 denote the complete graph on four vertices, and let G_0 be the graph obtained from the union of $k - 1$ vertex-disjoint copies of K_4 by adding a vertex connected to every other point. Let \mathcal{G}_0 be the family of all *subdivisions* of G_0 , i.e., the set of all graphs arising from G_0 by replacing its edges with independent paths. Clearly, \mathcal{G}_0 is a countable subfamily of \mathcal{G}_k . On the other hand, it is easy to check that, if G is a graph containing two subgraphs isomorphic to distinct elements of \mathcal{G}_0 , then $G \notin \mathcal{G}_k$. Thus, $c(\mathcal{G}_k) \geq |\mathcal{G}_0| = \omega$, completing the proof.

The analogous result for countable graphs containing no k edge-disjoint cycles can be established by a similar argument.

Theorem. *Let $1 < k < \omega$, and let \mathcal{G}'_k be the class of all countable graphs containing no k edge-disjoint cycles. Then $c(\mathcal{G}'_k) = \omega$.*

To see that $c(\mathcal{G}'_k) \geq \omega$, we can repeat the above argument with the only difference that now G_0 has to be defined as the graph obtained from the union of $k - 1$ vertex-disjoint *triangles* by adding a vertex connected to every other point. The minor modifications in the other part of the proof are left to the reader.

References

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