Scientific Computing, Courant Institute, Fall 2018

http://www.math.nyu.edu/faculty/goodman/teaching/ScientificComputing2018/ScientificComputing.html Always check the classes message board before doing any work on the assignment.

Assignment 3

Corrections: [none yet]

1. Suppose A(s) is a differentiable family of positive definite symmetric matrices. Suppose that A has a Cholesky factorization $A(s) = L(s)L^t(s)$. Here, L is lower triangular. Find a formula for

$$\dot{L} = \frac{d}{ds}L(s)$$

as a function of L, L^{-1} and \dot{A} . It is a useful fact that the inverse of a lower triangular matrix is lower triangular.

2. Suppose A is an $n \times n$ matrix. Show that the power series

$$e^{tA} = \sum_{0}^{\infty} \frac{1}{n!} t^n A^n \; .$$

converges in any matrix norm. Hint: Never mind the first N terms. What matters is the "tails", the terms $M_n = \frac{1}{n!}t^n A^n$ for $n \ge N$. Choose N so that $N \ge 2|t| ||A||$. Then

$$M_n = \frac{1}{n} t A M_{n-1}$$

Therefore, if $n \ge N$,

$$||M_n|| \le \frac{1}{N} |t| ||A|| ||M_{n-1}|| \le \frac{1}{2} ||M_{n-1}||$$

Therefore, if n > N, we have $||M_n|| \le \frac{1}{2^{n-N}} ||M_N||$. The tail sum satisfies

$$\left\|\sum_{n=N}^{\infty} M_n\right\| \le \left(\sum_{n\ge N} \frac{1}{2^{n-N}}\right) \|M_N\| \le 2 \|M_N\| < \infty.$$

The tail converges absolutely so the whole sum does too.

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- 3. (Properties of the matrix exponential)
 - (a) Show that the matrix $S(t) = e^{tA}$ satisfies a matrix differential equation and commutes with A:

$$\frac{d}{dt}S(t) = AS(t) = S(t)A \; .$$

- (b) Show that $S(t) \to I$ as $t \to 0$.
- (c) Show that $S(t_1)S(t_2) = S(t_1 + t_2)$. Hint: You can do this using part (d) below, or you can do it directly. For the direct argument, define $p = \frac{t_1}{t_1+t_2}$ and $q = \frac{t_2}{t_1+t_2}$ and write the formula $(p+q)^n = 1$ as

$$\sum_{j+k=n} \binom{n}{k} p^j q^k = \sum_{j+k=n} \frac{n!}{j!k!} p^j q^k = 1.$$

Then write

$$S(t_1)S(t_2) = \sum_j \sum_k \frac{1}{j!} \frac{1}{k!} t_1^j t_2^k A^{j+k} = \sum_n \left(\sum_{j+k=n} \frac{n!}{j!k!} p^j q^k \right) t^n A^n .$$

- (d) Suppose that $A = R\Lambda L$, where RL = I and Λ is the diagonal matrix of eigenvalues. Allow the possibility that R, L, and Λ are complex. Show that $S(t) = Re^{t\Lambda}L$. Show that $e^{t\Lambda}$ is diagonal and find formulas for the numbers on the diagonal.
- (e) The Python package linalg has routines for computing the eigenvalues and eigenvectors of a non-symmetric matrix. Use this to write a Python module that evaluates S(t) given A and t. Test your module by calling it for the case $A = \begin{pmatrix} 0 & 1 \\ -1 & 0 \end{pmatrix}$. In this case $S(t) = \begin{pmatrix} \cos(t) & \sin(t) \\ -\sin(t) & \cos(t) \end{pmatrix}$. Warning: because of roundoff in complex arithmetic, the computed result might not be real. The imaginary

part of the computed S(t) is all error.

- 4. Another way to compute the matrix exponential is to compute $S(\Delta t)$ using a few terms of the Taylor series, then take powers to apply the formula $S(n\Delta t) = S(\Delta t)^n$. Experiment with k terms of the Taylor series. Then choose n depending on t, and then $\Delta t = t/n$ so that k terms are enough to get $S(\Delta t)$ ten significant digits.
 - (a) Compute the Taylor series numerically using Horner's rule which is

$$x_0 I + x_1 A + \dots + x_{k-1} A^{k-1} + x_k A^k$$

= $x_0 I + x_1 A + \dots + x_{k-1} A^{k-1} \left(I + \frac{x_k}{x_{k-1}} A \right)$
= $x_0 I + x_1 A + \dots + x_{k-2} A^{k-2} \left(I + \frac{x_{k-1}}{x_{k-2}} A \left(I + \frac{x_k}{x_{k-1}} A \right) \right)$
= \dots
= $x_0 \left(I + \frac{x_1}{x_0} A \left(I + \frac{x_2}{x_1} A \left(\dots \right) \right) \right)$

This gives a way to calculate the k term sum using k matrix multiplications. Doing it directly takes $O(k^2)$ multiplications.

(b) Compute a power B^n in the following way. First compute $B_1 = B^2$, then $B_2 = B_1^2 = B^4$, then $B_3 = B_2^2 = B^{2^3}$, etc. Then write $n = n_0 + 2n_1 + 4n_2 + \cdots$, where $n_j = 0$ of $n_j = 1$ (expand n base 2) and compute

$$B^n = \prod_{n_j=1} B_j \; .$$

Show that this computes B^n in $O(\log(n))$ matrix multiplications.

- (c) Test your program on the A from problem 3.
- 5. Consider the differential equation system

$$\begin{aligned} \frac{dx_k}{dt} &= p(x_{k+1} - x_k) + q(x_{k-1} - x_k) & \text{if } 1 < k < n \\ \frac{dx_1}{dt} &= p(x_2 - x_1) \\ \frac{dx_n}{dt} &= q(x_{n-1} - x_n) . \end{aligned}$$

Here, x_k could represent the number of particles at spot k, and p represents the rate at which particles move to the left (from k + 1 to k or from k to k-1) and q represents the rate at which they move to the right. Particles cannot go from k = 1 to k = 0 or from k = n to k = n + 1. Try your codes from Problem 4 and Problem 5 on this system. If you explore p > q and nrelatively large, you should find that the eigenvalue method breaks down because the eigenvalue/eigenvector calculation is too ill conditioned. You will need code to create the matrix A. Note (check) that $M = \sum_k x_k(t)$ is a constant. Start with $x_1(n) = 1$, and $x_k(0) = 0$ for k < n. The solution has $x(t) \to x_*$ as $t \to \infty$. You can find x_* by setting $Ax_* = 0$ and $\sum_k x_k = 1$. Here are some things to explore as time and interest permits

- Find values of n, p > 0, q > 0, and t so that the eigenvalue method breaks down.
- The eigenvalues of A are all real. Monitor the imaginary parts of the computed eigenvalues.
- The problem of finding e^{tA} for this problem is well conditioned. The eigenvalue method is an unstable algorithm.
- 6. Write Python code to compute the matrix product AB when A and B are $n \times n$ matrices. You can initialize the matrices how you want because the values are unimportant. Do it in three ways:
 - Scalar loop, as

```
for i in range(n):
for j in range(n):
    c[i,j] = 0.
    for k in range(n):
        c[i,j] += a[i,k]*b[k,j]
```

• Hand coded vector loops, as (names and syntax wrong, please fix)

c[i,j] = np.sum(a[i,:]*b[:,j])

• The numpy matrix multiply routine, as

c = np.linalg.matrixMultiply(a, b)

Learn how to time sections of Python code. Print the times for the three methods as a function of n for an interesting range of n. Comment on the differences.